Sequential Investment and Online Prediction Tengyuan Liang¹ DLA Lecture 4: No Regret

Readings: Cesa-Bianchi and Lugosi², Chapter 9 and 10.

Contents

- 1 Sequential Portfolio Selection 1
- 1.1 Minimax Wealth Ratio 3
- 2 Digression: Online Prediction with Log Loss 4
- 2.1 Connection to Maximum Likelihood Estimation
- 2.2 Connection to Bayesian Algorithm
- 3 Reduction: Sequential Investment to Online Prediction 7

5

7

5

- 4 Universal Portfolio and Bayesian Mixtures
- 5 Exponentiated Gradient (EG) Portfolio 10
- 6 Summary 11

1 Sequential Portfolio Selection

Consider the problem of sequential investment. A market consists of m stocks in which, in each trading period $t \in 1, 2...n$, the price of a stock may vary in an arbitrary way. An investor seeks to sequentially allocate the portfolio in a trading period of n days.

We deliberately avoid any statistical assumptions about the nature of the stock market, and evaluate the investor's wealth relative to the performance achieved by the best strategy in a class of reference investment strategies, or the so-called experts. We assume no transaction costs.

Let's fix the notations.

- Market information: x = (x₁,..., x_m)^T ∈ ℝ^m_{≥0} denotes the ratio of closing to opening price of the i-*th* stock in that period. Since we consider a trading period t = 1, 2, ..., n, we stack the market information at time t to be the matrix x^t := [x₁,..., x_t] ∈ ℝ^{m×t}. Again, the *i*-th component of x_t is denoted as x_{i,t}, is the factor by which the wealth invested in stock *i* increases in the period t.
- Investment strategy: an investment vector **q** = (q₁,..., q_m) ∈ Δ_m the probability simplex denotes an allocation of the total wealth

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² Nicolò Cesa-Bianchi and Gábor Lugosi. *Prediction, Learning, and Games.* Cambridge University Press, 2006 to each stock. If invest a unit amount to the market vector **x** with investment strategy **q**, by the end of the trading, the wealth grows to $\sum_{i=1}^{m} q_i x_i$.

A sequential investment strategy denotes a mapping from past market information to a probability distribution over the stocks, namely $\mathbf{Q}_t : \mathbb{R}_{>0}^{m \times (t-1)} \to \Delta_m$, which specifies

$$\mathbf{x}^{t-1} \mapsto \mathbf{Q}_t(\mathbf{x}^{t-1}) \in \Delta_m, \ t = 1, 2, \dots, n \tag{1.1}$$

We denote $Q_{i,t}(\mathbf{x}^{t-1})$ to denote the portation of total wealth invested to stock *i* in period *t*, based on market information collected in the past t-1 periods. We call the sequential investment strategy (for periods n) $\mathbf{Q} = (\mathbf{Q}_1, \mathbf{Q}_2, \dots, \mathbf{Q}_n)$ to be the collection of such maps.

• Wealth factor: given the martket information *x*^{*n*} and a sequantial investment stragegy **Q**, compute the wealth factor

$$S_n(\mathbf{Q}, \mathbf{x}^n) := \prod_{t=1}^n \sum_{i=1}^m Q_{i,t}(\mathbf{x}^{t-1}) x_{i,t}$$
(1.2)

Now, we consider two sets of simple strategies.

Example (Buy-and-hold stragegy). A buy-and-hold stragegy refers to a no-trading strategy indexed by $q \in \Delta_m$, where one distributes wealth according to q and hold it for a period of n,

$$S^{\text{BnH}}(\mathbf{q}, \mathbf{x}^n) := \sum_{i=1}^m \mathbf{q}_i \prod_{t=1}^n x_{i,t}$$
 (1.3)

Example (Constantly rebalancing strategy). A constantly rebalancing strategy refers to a time-homogeneous strategy that $\mathbf{Q}_t(\mathbf{x}^{t-1}) := \mathbf{q}$ regardless of time t and the past market behavior \mathbf{x}^{t-1} . Note there we trade each period, thus bearing the name constantly rebalancing (as opposed to the buy-and-hold strategy, where one does not trade at all).

$$S_n^{\text{ReB}}(\mathbf{q}, \mathbf{x}^n) := \prod_{t=1}^n \sum_{i=1}^m \mathbf{q}_i x_{i,t} .$$
 (1.4)

We call the constantly rebalancing strategy

$$\mathcal{B}^{\text{ReB}} := \{ (\mathbf{q}, \mathbf{q}, \dots, \mathbf{q}), \mathbf{q} \in \Delta_m \} .$$
(1.5)

To better constrast the buy-and-hold strategy and the constantly rebalacing strategy, consider a simple case where there are two stocks,

$$\mathbf{x}^{n} := \begin{bmatrix} 1 & 1 & 1 & 1 & \dots & 1 & 1 \\ 0.5 & 2 & 0.5 & 2 & \dots & 0.5 & 2 \end{bmatrix}$$
(1.6)

Then for n even

$$S^{ ext{BnH}}(\mathbf{q},\mathbf{x}^n) = 1, orall q$$

 $S^{ ext{ReB}}_n((1/2,1/2),\mathbf{x}^n) = (rac{9}{8})^{n/2}$

1.1 Minimax Wealth Ratio

Recall the goal of the investor is to compete with a certain class of investment strategies Q, regardless of the market behavior. Such classes might include, actively managed portfolios built based on these *m* stocks, or all constantly rebalancing strategies.

Definition 1 (Minimax wealth ratio). *The worst-case logarithmic wealth ratio* of a strategy **P** relative to a class Q is defined as

$$W_n(\mathbf{P}, \mathcal{Q}) := \sup_{\mathbf{x}^n} \sup_{\mathbf{Q} \in \mathcal{Q}} \log \frac{S_n(\mathbf{Q}, \mathbf{x}^n)}{S_n(\mathbf{P}, \mathbf{x}^n)}$$
(1.7)

and define the minimax wealth ratio as

$$W_n(\mathcal{Q}) := \inf_{\mathbf{P}} W_n(\mathbf{P}, \mathcal{Q}) .$$
(1.8)

A few remarks follow

• For a good class of reference strategies and good market conditions, we can expect the wealth factor grow exponentially

$$\sup_{\mathbf{Q}\in\mathcal{Q}} \log S_n(\mathbf{Q}, \mathbf{x}^n) = r \cdot n \tag{1.9}$$

For instance, the S&P500 has a growth rate of $\exp(\log(1.095)n)$, where r = 0.04.

• If a strategy **P** satisfies

$$W_n(\mathbf{P}, \mathcal{Q}) = o(n)$$

then it implies that **P** can compete with the class \mathcal{Q}

$$\log S_n(\mathbf{P}, \mathbf{x}^n) \geq \underbrace{\log S_n(\mathbf{Q}, \mathbf{x}^n)}_{=r \cdot n} - o(n) .$$

The name no-regret alludes to the above property.

Example (Compete with *N* portfolio managers/experts). Consider Q to be the class of strategies presented by *N* portfolio managers,

$$\mathcal{Q} := \{\mathbf{Q}^{(1)}, \dots, \mathbf{Q}^{(N)}\}\$$

Then then simple average strategy **P** which divide the initial wealth in 1/N, 1/N, ..., 1/N and then invest on each expert using the buy-and-hold strategy, we have

$$S_n(\mathbf{P}, \mathbf{x}^n) = \frac{1}{N} \sum_{j=1}^N S_n(\mathbf{Q}^{(j)}, \mathbf{x}^n)$$

and thus

$$W_n(\mathbf{P}, \mathcal{Q}) \leq \log(N)$$

2 Digression: Online Prediction with Log Loss

Consider the online probability assignment question, where there are m items $\mathcal{Y} = \{1, 2, ..., m\}$ and n-periods. One wishes to maximum the likelihood of the sequence by assigning the probabilities.

- **Sequences**: A sequence of numbers is revealed sequentially (y_1, y_2, \ldots, y_n) , where each $y_t \in \mathcal{Y}$, finite *m*-items. We denote $y^t := (y_1, y_2, \ldots, y_t)$.
- **Experts**: An expert $q = (\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_n)$ is a sequence of functions $\mathbf{q}_t : \mathcal{Y}^{t-1} \to \Delta_m$

$$y^{t-1} \mapsto \mathbf{q}_t(y^{t-1}) \in \Delta_m \tag{2.1}$$

is a probability assignment vector, with each component

$$\mathbf{q}_t(y^{t-1}) := [q_t(1|y^{t-1}), q_t(2|y^{t-1}), \dots, q_t(m|y^{t-1})].$$

We denote the class of experts as \mathcal{E} .

- Forecaster: **p** = (**p**₁, **p**₂, ..., **p**_n) is a sequence of probability vectors similar to the definition in experts.
- **Likelihood**: One more notation, given a forecaster *p* (or similarly for an expert *q*), we define the likelihood for the sequence *y*^{*n*}

$$p_n(y^n) := \prod_{t=1}^n p_t(y_t | y^{t-1}) .$$
(2.2)

Note here the goal is to maximize the log-likelihood in comparison to a class of Experts \mathcal{E}

$$M_n(\mathbf{p}, \mathcal{E}) := \max_{y^n \in \mathcal{Y}^n} \left\{ \sup_{\mathbf{q} \in \mathcal{E}} \log q_n(y^n) - \log p_n(y^n) \right\}$$
(2.3)

2.1 Connection to Maximum Likelihood Estimation

Definition 2 (Minimax Optimal Forecaster). *Define the minimax regret for the class* \mathcal{E} *as*

$$M_n(\mathcal{E}) := \inf_{\mathbf{p}} M_n(\mathbf{p}, \mathcal{E})$$

The infimum turns out to be attained by the normalized maximum likelihood probability distribution (MLE)

$$p_n^{\star}(y^n) := \frac{\sup_{\mathbf{q}\in\mathcal{E}} q_n(y^n)}{\sum_{z^n\in\mathcal{Y}^n} \sup_{\mathbf{q}\in\mathcal{E}} q_n(z^n)}$$

and the sequential forecaster is thus defined as

$$p_t^{\star}(i|y^{t-1}) := \frac{p_t^{\star}(y^{t-1}i)}{p_{t-1}^{\star}(y^{t-1})} .$$
(2.4)

Note that this minimax optimal forecaster is a computation burden to calculate, even for m = 2.

Theorem 1 (Regret for Minimax Optimal Forecaster). *The minimax regret is achieved by the minimax optimal forecaster, and*

$$M_n(\mathcal{E}) = \log\left(\sum_{y^n \in \mathcal{Y}^n} \sup_{\mathbf{q} \in \mathcal{E}} q_n(z^n)\right)$$
(2.5)

How to deal with the computation then? The trick is to swap the "sup" by " \int " over a distribution of q_t denoted by $\mu \in \mathcal{P}(\mathcal{E})$.

$$p_n^{\mu}(y^n) := \frac{\int_{q \in \mathcal{E}} q_n(y^n) \, d\mu(q)}{\sum_{z^n \in \mathcal{Y}^n} \int_{q \in \mathcal{E}} q_n(z^n) \, d\mu(q)}$$
$$= \frac{\int_{q \in \mathcal{E}} q_n(y^n) \, d\mu(q)}{\int_{q \in \mathcal{E}} (\sum_{z^n \in \mathcal{Y}^n} q_n(z^n)) \, d\mu(q)} \quad \text{Fubini's theorem}$$
$$= \int_{q \in \mathcal{E}} q_n(y^n) \, d\mu(q)$$

It turns out this integration swapping supermum idea also has a nice Bayesian interpretation, which we will discuss now.

2.2 Connection to Bayesian Algorithm

Definition 3 (Bayesian Mixture Forecaster). *Given a prior distribution, which is a mixture of experts in* \mathcal{E} *, denoted by* μ *, define the following Bayesian mixture forecaster*

$$p_n^{\mu}(y^n) := \int_{q \in \mathcal{E}} q_n(y^n) \,\mathrm{d}\mu(q) \tag{2.6}$$

and correspondingly

$$p_t^{\mu}(i|y^{t-1}) := \frac{\int_{q \in \mathcal{E}} q_t(y^{t-1}i) \, \mathrm{d}\mu(q)}{\int_{q \in \mathcal{E}} q_{t-1}(y^{t-1}) \, \mathrm{d}\mu(q)} = \Pr_{\text{posterior}}(i|y^{t-1}) \tag{2.7}$$

where the posterior probability is calculated based on the prior $q \sim \mu$, and then $y^t \sim q_n(y^n)$.

Theorem 2 (Regret for Bayesian Mixture Forecaster). *The minimax regret is achieved by the minimax optimal forecaster, and*

$$M_n(p^{\mu}, \mathcal{E}) = \sup_{y^n \in \mathcal{Y}^n} \log \frac{\sup_{q \in \mathcal{E}} q_n(y^n)}{\int_{q \in \mathcal{E}} q_n(y^n) \, \mathrm{d}\mu(q)}$$
(2.8)

Definition 4 (Constant Experts). If the experts are time-homogeneous, namely, $q_t(i|y^{t-1}) \equiv q(i), \forall t$, we call them constant experts. For this problem, we define the constant experts class as

$$\mathcal{E}^{\mathbf{h}} := \{(\underbrace{\mathbf{q}, \mathbf{q}, \dots, \mathbf{q}}_{n}) \mid \mathbf{q} \in \Delta_{m}\}$$
(2.9)

Theorem 3 (Regret Bounds: Krichevsky-Trofimov mixture forecaster vs. minimax optimal forecaster). *Consider the case of constant experts, then the minimax optimal forecaster achieves the regret*

$$M_n(\mathcal{E}^{\rm h}) = \frac{m-1}{2}\log\frac{n}{2\pi} + \log\frac{\Gamma(1/2)^m}{\Gamma(m/2)} + o_n(1) .$$
 (2.10)

Pick $\mu(\mathbf{q})$ *where* $\mathbf{q} \in \Delta_m$ *and* μ *being the Dirichlet prior*

$$d\mu(\mathbf{q}) = \frac{\Gamma(m/2)}{\Gamma^m(1/2)} \prod_{i=1}^m \frac{1}{\sqrt{q(i)}} \, d\mathbf{q}$$
(2.11)

and consider the Bayesian mixture forecaster defined in (2.6)

$$p_n^{\mu}(y^n) := \int_{\Delta_m} \prod_{t=1}^n q(y_t) \, \mathrm{d}\mu(\mathbf{q})$$
 (2.12)

Then the corresponding regret bound holds

$$M_n(p^{\mu}, \mathcal{E}^{\mathbf{h}}) := \sup_{y^n \in \mathcal{Y}^n} \sup_{q \in \mathcal{E}^{\mathbf{h}}} \left\{ \log q_n(y^n) - \log p_n^{\mu}(y^n) \right\}$$
(2.13)

$$\leq \frac{m-1}{2}\log\frac{n}{2\pi} + \log\frac{\Gamma(1/2)^m}{\Gamma(m/2)} + \frac{m-1}{2}\log 2 + o_n(1) .$$
 (2.14)

A few remarks follow

• The Krichevsky-Trofimov mixture forecaster exceeds the minimax optimal bound just by a constant factor (independent of *n*). In particular,

$$M_n(p^{\mu}, \mathcal{E}^{\mathbf{h}}) \leq \underbrace{\inf_p M_n(p^{\mu}, \mathcal{E}^{\mathbf{h}})}_{\asymp \frac{m-1}{2}\log n} + \frac{m-1}{2}\log 2$$
(2.15)

The proof is based on Gamma functions and Stirling approximations.

• The Krichevsky-Trofimov mixture may be easily calculated by a smoothed version of empirical frequencies

$$p_t^{\mu}(i|y^{t-1}) = \frac{t_i + 1/2}{t - 1 + m/2}$$
(2.16)

where t_i denotes the number of occurrences of i in y^{t-1} with $\sum_{i=1}^{m} t_i = t - 1$.

3 Reduction: Sequential Investment to Online Prediction

In this section, we will show a reduction to relate the two problems: sequential investment, and online probability assignment.

The idea is simple based on two steps:

- For a minimax sequential investment problem, restricted **only a strict subset of market conditions** (called Kelly market vectors), we reduce to a minimax online probability assignment problem. This will show the minimax sequential investment problem is strictly harder than the online prediction problem.
- Given any algorithm *p* that solves the online probability assignment problem, we induce a corresponding sequential investment algorithm **P** for **any market conditions**. This step hopes to upper bound the regret of **P** for the sequential investment problem by the regret of *p* for the online learning problem.

4 Universal Portfolio and Bayesian Mixtures

Define the initial wealth $S_0 \equiv 1$ and define the wealth at *t*

$$S_t^{\text{ReB}}(\mathbf{q}, \mathbf{x}^t) := \prod_{s=1}^t \left(\sum_{i=1}^m \mathbf{q}_i x_{i,s}\right)$$
(4.1)

Define **the universal portfolio strategy** $\mathbf{P}_{\mathrm{UP}}^{\mu}$ induced by μ

$$P_{i,t}^{\mathrm{UP}}(\mathbf{x}^{t-1}) := \frac{\int_{\Delta_m} \mathbf{q}_i \cdot S_{t-1}^{\mathrm{ReB}}(\mathbf{q}, \mathbf{x}^{t-1}) \,\mathrm{d}\mu(\mathbf{q})}{\int_{\Delta_m} S_{t-1}^{\mathrm{ReB}}(\mathbf{q}, \mathbf{x}^{t-1}) \,\mathrm{d}\mu(\mathbf{q})}$$
(4.2)

In any case, the universal portfolio is a weighted average of the strategies in \mathcal{B}^{ReB} , weighted by their past performance.

TL: Leave it as homework. It is not hard to show.

Observe that

$$S_{n}(\mathbf{P}^{\mu}, \mathbf{x}^{n}) = \prod_{t=1}^{n} \sum_{i=1}^{m} P_{i,t}(\mathbf{x}^{t-1}) x_{i,t}$$

$$= \prod_{t=1}^{n} \frac{\int_{\Delta_{m}} \sum_{i=1}^{m} \mathbf{q}_{i} x_{i,t} \cdot S_{t-1}^{\text{ReB}}(\mathbf{q}, \mathbf{x}^{t-1}) d\mu(\mathbf{q})}{\int_{\Delta_{m}} S_{t-1}^{\text{ReB}}(\mathbf{q}, \mathbf{x}^{t-1}) d\mu(\mathbf{q})}$$

$$= \prod_{t=1}^{n} \frac{\int_{\Delta_{m}} S_{t}^{\text{ReB}}(\mathbf{q}, \mathbf{x}^{t}) d\mu(\mathbf{q})}{\int_{\Delta_{m}} S_{t-1}^{\text{ReB}}(\mathbf{q}, \mathbf{x}^{t-1}) d\mu(\mathbf{q})}$$

$$= \int_{\Delta_{m}} S_{n}^{\text{ReB}}(\mathbf{q}, \mathbf{x}^{n}) d\mu(\mathbf{q}) = \int_{\Delta_{m}} \prod_{t=1}^{n} (\sum_{i=1}^{m} \mathbf{q}_{i} x_{i,t}) d\mu(\mathbf{q})$$

$$= \int_{\Delta_{m}} \sum_{y_{n} \in \mathcal{Y}^{n}} \prod_{t=1}^{n} \mathbf{q}_{y_{t}} x_{y_{t},t} d\mu(\mathbf{q})$$

$$= \sum_{y_{n} \in \mathcal{Y}^{n}} (\prod_{t=1}^{n} x_{y_{t},t}) \underbrace{\int_{\Delta_{m}} \prod_{t=1}^{n} \mathbf{q}_{y_{t}} d\mu(\mathbf{q})}_{p_{n}^{\mu}(y^{n})}$$

which is governed by the online prediction problem.

Algorithm: universal portfolio

Input: a sequence of market vectors X_{i,t}, i from m assets, t from n periods # Output: Krichevsky-Trofimov mixture forecaster to build universal portfolio def log_wealth_factor_rebalance(Q, X):

```
# m, n = np.shape(X)
# Nsim, m = np.shape(Q)
lnW = np.log(np.matmul(Q, X))
lnWcum = np.cumsum(lnW, axis = 1)
return lnWcum
```

```
def universal_portfolio(X, Nsim = 1e5):
    m, n = np.shape(X)
    alpha = 0.5 * np.ones(m)
    Q = np.random.dirichlet(alpha, int(Nsim))
```

lnWcum = log_wealth_factor_rebalance(Q, X)
P_un_normalized = np.matmul(np.transpose(Q), np.exp(lnWcum-np.log(Nsim)))
norm_Const = np.sum(np.exp(lnWcum-np.log(Nsim)), axis = o, keepdims=True)
return P_un_normalized/norm_Const

Theorem 4 (Performance of Universal Portfolio). *The universal portfolio algorithm* $\mathbf{P}^{\mu}_{\text{UP}}$ *above satisfies the following regret guarantee over the class of constantly rebalancing strategies defined in* (1.5),

 $\sup_{\mathbf{x}^n} \sup_{\mathbf{Q} \in \mathcal{B}^{\text{ReB}}} \log S_n(\mathbf{Q}, \mathbf{x}^n) - \log S_n(\mathbf{P}_{UP}^{\mu}, \mathbf{x}^n) \leq \frac{m-1}{2} \log \frac{n}{2\pi} + \log \frac{\Gamma(1/2)^m}{\Gamma(m/2)} + \frac{m-1}{2} \log 2 + o_n(1)$

Proof. The proof hinges on a reduction to the online probability assignment problem, as hinted above.

Note that for any $\mathbf{Q} \in \mathcal{B}^{\text{ReB}}$, it naturally induced an online probability assignment problem, and therefore: Given \mathbf{x}^n , let \mathbf{q}^\dagger be the maximizer of $\sup_{\mathbf{Q} \in \mathcal{B}^{\text{ReB}}} \log S_n(\mathbf{Q}, \mathbf{x}^n)$, then

$$\begin{split} \sup_{\mathbf{Q}\in\mathcal{B}^{\text{ReB}}} \log S_n(\mathbf{Q},\mathbf{x}^n) &- \log S_n(\mathbf{P}_{UP}^{\mu},\mathbf{x}^n) \\ &= \log S_n(\mathbf{Q}^{\dagger},\mathbf{x}^n) - \log S_n(\mathbf{P}_{UP}^{\mu},\mathbf{x}^n) \\ &= \log \frac{\sum_{y_n\in\mathcal{Y}^n}(\prod_{t=1}^n x_{y_t,t})\prod_{t=1}^n \mathbf{q}_{y_t}^{\dagger}}{\sum_{y_n\in\mathcal{Y}^n}(\prod_{t=1}^n x_{y_t,t})\int_{\Delta_m}\prod_{t=1}^n \mathbf{q}_{y_t} \, \mathrm{d}\mu(\mathbf{q})} \\ &= \log \frac{\sum_{y_n\in\mathcal{Y}^n}(\prod_{t=1}^n x_{y_t,t}) \cdot q_n^{\dagger}(y^n)}{\sum_{y_n\in\mathcal{Y}^n}(\prod_{t=1}^n x_{y_t,t}) \cdot p_n^{\mu}(y^n)} \end{split}$$

Denote $w(y^n, \mathbf{x}^n) := \prod_{t=1}^n x_{y_t,t} \in \mathbb{R}_{\geq 0}$ to be some non-negative weights, then the above

$$\log \frac{\sum_{y_n \in \mathcal{Y}^n} w(y^n, \mathbf{x}^n) \cdot q_n^{\dagger}(y^n)}{\sum_{y_n \in \mathcal{Y}^n} w(y^n, \mathbf{x}^n) \cdot p_n^{\mu}(y^n)} \\ \leq \sup_{y_n: w(y^n, \mathbf{x}^n) > 0} \log \frac{q_n^{\dagger}(y^n)}{p_n^{\mu}(y^n)} \\ \leq \sup_{y_n: w(y^n, \mathbf{x}^n) > 0} \sup_{\mathbf{q} \in \mathcal{E}^h} \log \frac{q_n(y^n)}{p_n^{\mu}(y^n)} \\ \leq \sup_{y_n \in \mathcal{Y}^n} \sup_{\mathbf{q} \in \mathcal{E}^h} \log \frac{q_n(y^n)}{p_n^{\mu}(y^n)}$$

Put things together, we have derived that the upper bound of the induced universal portfolio algorithm is bounded by that of the online Krichevsky-Trofimov mixture forecaster

 $\sup_{\mathbf{x}^{n}} \sup_{\mathbf{Q} \in \mathcal{B}^{\text{ReB}}} \log S_{n}(\mathbf{Q}, \mathbf{x}^{n}) - \log S_{n}(\mathbf{P}_{UP}^{\mu}, \mathbf{x}^{n}) \leq \sup_{y_{n} \in \mathcal{Y}^{n}} \sup_{\mathbf{q} \in \mathcal{E}^{\text{h}}} \log q_{n}(y^{n}) - \log p_{n}^{\mu}(y^{n})$ (4.3)

and by Theorem 3, we have the RHS is bounded by the desired quantity. $\hfill \Box$

Remark. It turns out the equivalence is stronger; one can show that the sequential portfolio question is precisely as hard as the online probability assignment question, in the minimax sense.

$$W_n(\mathcal{B}^{\text{ReB}}) = \inf_{\mathbf{P}} \sup_{\mathbf{x}^n} \left\{ \sup_{\mathbf{Q} \in \mathcal{B}^{\text{ReB}}} \log S_n(\mathbf{Q}, \mathbf{x}^n) - \log S_n(\mathbf{P}, \mathbf{x}^n) \right\}$$
$$= \inf_{\mathbf{P}} \sup_{y_n \in \mathcal{Y}^n} \left\{ \sup_{\mathbf{q} \in \mathcal{E}^{\text{h}}} \log q_n(y^n) - \log p_n(y^n) \right\}$$
$$= M_n(\mathcal{E}^{\text{h}})$$

The \geq can be shown by considering only Kelly vectors ³ as market information. See Theorem 10.1 in ⁴.

5 Exponentiated Gradient (EG) Portfolio

The **the universal portfolio strategy** \mathbf{P}_{UP}^{μ} still involves multi-dimensional integrals in the space $\mathbf{q} \in \Delta_m$, and one has to calculate O(mn) such *m*-dimensional simplex integrals. Though an online algorithm, the procedure is computationally heavy and relies on Monte Carlo simulations to calculate.

In this section, we consider a linearized version that is much faster to compute, based on exponentiated gradient (EG) principle. Unfortunately, unlike the universal portfolio $\mathbf{P}_{\text{UP}}^{\mu}$ which has $\log(n)$ regret, the EG strategy has \sqrt{n} regret.

Consider the exponentiated gradient (EG) strategy $\mathbf{P}_{\text{EG}}^{\eta}$, an online algorithm. Define $P_{i,1}^{\text{EG}} = 1/m, \forall i \in [m]$ and update

$$P_{i,t}^{\text{EG}} = \frac{P_{i,t-1}^{\text{EG}} \exp\left(\frac{\eta x_{i,t-1}}{\langle \mathbf{P}_{t-1}^{\text{EG}}, \mathbf{x}_{t-1} \rangle}\right)}{\sum_{j=1}^{m} P_{j,t-1}^{\text{EG}} \exp\left(\frac{\eta x_{j,t-1}}{\langle \mathbf{P}_{t-1}^{\text{EG}}, \mathbf{x}_{t-1} \rangle}\right)}$$
(5.1)

Here, no integration is needed and the algorithm can be implemented exactly fast.

Motivation behind: Write out the total wealth for strategy $\mathbf{P} = (\mathbf{P}_t)_{t=1}^n$

$$-\log S_n(\mathbf{P}, \mathbf{x}^n) = \sum_{t=1}^n -\log \langle \mathbf{P}_t, \mathbf{x}_t \rangle$$
(5.2)

Define $\ell_t(\mathbf{P}) = -\log \langle \mathbf{P}, \mathbf{x}_t \rangle$ is a convex function in \mathbf{P} , and thus one can use the online EG algorithm,

$$P_{i,t} \propto P_{i,t-1} \exp(-\eta \nabla \ell_{t-1} (\mathbf{P}_{t-1})_i)$$
(5.3)

$$\propto P_{i,t-1} \exp\left(\frac{\eta x_{i,t-1}}{\langle \mathbf{P}_{t-1}, \mathbf{x}_{t-1} \rangle}\right)$$
(5.4)

therefore we derive

$$P_{i,t} = \frac{P_{i,t-1} \exp\left(\frac{\eta x_{i,t-1}}{\langle \mathbf{P}_{t-1}, \mathbf{x}_{t-1} \rangle}\right)}{\sum_{j=1}^{m} P_{j,t-1} \exp\left(\frac{\eta x_{j,t-1}}{\langle \mathbf{P}_{t-1}, \mathbf{x}_{t-1} \rangle}\right)}$$
(5.5)

Algorithm: exponentiated gradient portfolio def expGrad_portfolio(X, eta = 0.1):

m, n = np.shape(X)P = np.ones((m, n+1))/mfor t in range(n): ³ John L Kelly. A new interpretation of information rate. *the bell system technical journal*, 35(4):917–926, 1956
⁴ Nicolò Cesa-Bianchi and Gábor Lugosi. *Prediction, Learning, and Games*. Cambridge University Press, 2006 w = np.sum(P[:, t]*X[:, t])
P_un_normalized = P[:, t]*np.exp(eta*X[:, t]/w)
P[:, t+1] = P_un_normalized/np.sum(P_un_normalized)
return P

Theorem 5 (Performance of Exponentiated Gradient). Assume that the price relatives $x_{i,t}$ all fall between two positive constants c < C. Then EG portfolio with the $\eta = \frac{c}{C} \sqrt{\frac{8 \log m}{n}}$

$$\sup_{\mathbf{x}^{n} \in [c,C]^{m \times n}} \sup_{\mathbf{Q} \in \mathcal{B}^{\text{ReB}}} \log S_{n}(\mathbf{Q}, \mathbf{x}^{n}) - \log S_{n}(\mathbf{P}_{EG}, \mathbf{x}^{n})$$

$$\leq \frac{C}{c} \sqrt{\frac{n \log m}{2}}$$
(5.7)

The proof is a simple application of the optimistic mirror descent framework to analyze online problems with KL divergence as the Bregman divergence.

6 Summary

Note the difference in the wealth ratios

$$W_n(\mathbf{P}_{EG}, \mathcal{B}^{\text{ReB}}) \le \frac{C}{c} \sqrt{\frac{n \log m}{2}}$$
(6.1)

$$W_n(\mathbf{P}_{UP}, \mathcal{B}^{\text{ReB}}) \le \frac{m-1}{2} \log \frac{n}{2\pi}$$
(6.2)

Let's put all the algorithms shoulder-by-shoulder as a contrast

• The MLE forecaster effectively solves the following optimization problem

$$\min_{\mathbf{q}\in\Delta_m} \sum_{t=1}^n -\log\langle \mathbf{q}, \mathbf{x}_t \rangle \tag{6.3}$$

• The Dirichlet mixture forecaster (universal portfolio) solves the following optimization by sampling

$$\mathbf{q} \sim \exp(-F_n(\cdot)) \tag{6.4}$$

$$F_n(\mathbf{q}) := \sum_{t=1}^n -\log\langle \mathbf{q}, \mathbf{x}_t \rangle - \frac{1}{2} \langle 1, \log \mathbf{q} \rangle$$
(6.5)

• Tthe EG forecaster solves the following optimization iteratively using online mirror descent (or linearized the problem)

$$G_n(\mathbf{q}) := \sum_{t=1}^n -\log\langle \mathbf{q}, \mathbf{x}_t \rangle + \frac{1}{\eta} \langle \mathbf{q}, \log \mathbf{q} \rangle$$
(6.6)

namely

$$\mathbf{q}_{t} := \underset{\mathbf{q} \in \Delta_{m}}{\operatorname{arg\,min}} \langle \mathbf{q}_{\prime} - \frac{\mathbf{x}_{t-1}}{\langle \mathbf{q}_{t-1}, \mathbf{x}_{t} \rangle} \rangle + \frac{1}{\eta} \operatorname{KL}(\mathbf{q} \| \mathbf{q}_{t-1})$$
(6.7)

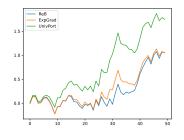


Figure 1: A simulation of the portfolios.

References

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