

# Reinforcement Learning II

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## DLA Lecture 6: Explore vs. Exploit

Bandit algorithms and details about exploration vs. exploitation. Readings: Hardt and Recht <sup>2</sup> Chapters 11 and 12, Lattimore and Szepesvari <sup>3</sup> Chapters 6, 7 and 11.

### Contents

1	Problem Setup	1
2	Stochastic Bandits	2
2.1	Explore-Then-Commit (ETC) Algorithm	2
2.2	Upper Confidence Bound (UCB) Algorithm	4
3	Adversarial Bandits	7
3.1	Exponential Weight for Exploration and Exploitation (EXP <sub>3</sub> ) Algorithm	7

### 1 Problem Setup

Multi-arm bandit <sup>4</sup> is a special class of sequential decision-making (SDM) problems we have considered before with no dynamic systems for how the state evolves: one simply decides based on the history the current action to take from a finite dictionary of arms, and then observes the reward. In other words, there is no states  $X$ , the reward can be represented by a vector of length  $k$ , and the action space is  $\mathcal{A}_k = \{1, 2, \dots, k\}$ .

**Definition 1** (Stochastic Bandit). Let the action space be  $\mathcal{A}_k = \{1, 2, \dots, k\}$ . Let the reward vectors  $\mathbf{r}_t \in \mathbb{R}^k, t = 1, 2, \dots, T$  be i.i.d. samples from an unknown rewards distribution. We denote the average reward for arm  $i$  as  $\mu(i) := \mathbb{E}[\mathbf{r}_t(i)], i = 1, 2, \dots, k$ .

At each time, the player takes an action  $u_t \in \mathcal{A}_k$  (based on the past information  $\{u_s, \mathbf{r}_s(u_s)\}_{s < t}$ ), then only observes the reward  $\mathbf{r}_t(u_t) \in \mathbb{R}$ . The player tries to maximize the cumulative reward

$$\max_{u_t} \mathbb{E} \left[ \sum_{t=1}^T \mathbf{r}_t(u_t) \right] = \max_{u_t} \mathbb{E} \left[ \sum_{t=1}^T \mu(u_t) \right].$$

Define the cumulative regret for an algorithm/policy  $\pi$ , where the actions  $u_t \sim \pi$

$$\mathcal{R}_T^{\text{sto}}(\pi) := T \cdot \max_{i \in \mathcal{A}_k} \mu(i) - \mathbb{E} \left[ \sum_{t=1}^T \mu(u_t) \right]$$

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<sup>2</sup> Moritz Hardt and Benjamin Recht. *Patterns, Predictions, and Actions: A Story about Machine Learning*. Princeton University Press, 2022

<sup>3</sup> Tor Lattimore and Csaba Szepesvári. *Bandit algorithms*. Cambridge University Press, 2020

<sup>4</sup> William R Thompson. On the likelihood that one unknown probability exceeds another in view of the evidence of two samples. *Biometrika*, 25(3-4): 285–294, 1933; and Herbert Robbins. Some aspects of the sequential design of experiments. *Bull. Amer. Math. Soc.*, 58(6):527–535, 1952

**Definition 2** (Adversarial Bandit). Let the action space be  $\mathcal{A}_k = \{1, 2, \dots, k\}$ . The reward vectors  $\mathbf{r}_t \in \mathbb{R}^k, t = 1, 2, \dots, T$  now can be arbitrary vectors.

At each time, the player takes an action  $u_t \in \mathcal{A}_k$  (based on the past information  $\{u_s, \mathbf{r}_s(u_s)\}_{s < t}$ ), then observes the reward  $\mathbf{r}_t(u_t) \in \mathbb{R}$ . The player tries to maximize the cumulative reward

$$\max_{u_t} \sum_{t=1}^T \mathbf{r}_t(u_t).$$

Define the cumulative regret for an algorithm/policy  $\pi$ , where the actions  $u_t \sim \pi$

$$\mathcal{R}_T^{\text{adv}}(\pi) := \max_{i \in \mathcal{A}_k} \sum_{t=1}^T \mathbf{r}_t(i) - \mathbb{E} \left[ \sum_{t=1}^T \mathbf{r}_t(u_t) \right]$$

## 2 Stochastic Bandits

### 2.1 Explore-Then-Commit (ETC) Algorithm

**Definition 3** (ETC Algorithm). *Explore-Then-Commit* (ETC) Algorithm specifies an exploration budget of time  $m \times k$ .

Define

$$N^t(i) := \sum_{s=1}^t \mathbf{1}_{u_s=i} \quad (2.1)$$

$$\hat{\mu}^t(i) := \frac{1}{N^t(i)} \sum_{s=1}^t \mathbf{1}_{u_s=i} \cdot \mathbf{r}_s(i) \quad (2.2)$$

which correspond to the number of times action  $i$  is taken up till time  $t$ , and the empirical estimate of the average reward for arm  $i$ .

The ETC algorithm  $\pi^{\text{ETC}}$  implements the following

1. Input an integer  $m$
2. In round  $t$ , choose action

$$u_t = \begin{cases} (t-1 \bmod k) + 1 & \text{if } t \leq mk \\ \arg \max_i \hat{\mu}^{mk}(i) & \text{if } t > mk \end{cases}$$

One more notation we need is the gap of the problem

$$\Delta_i := \max_{i'} \mu(i') - \mu_i, \quad i \in [k]. \quad (2.3)$$

**Theorem 1** (Regret for ETC Algorithm). Consider stochastic bandits with bounded rewards  $\mathbf{r}(i) \in [-1/\sqrt{2}, 1/\sqrt{2}]$ .

$$\mathcal{R}_T^{\text{sto}}(\pi^{\text{ETC}}) \leq \sum_{i=1}^k m \Delta_i + (T - mk) \Delta_i \exp\left(-\frac{m \Delta_i^2}{4}\right) \quad (2.4)$$

TL: The constants are for convenience only.

A few remarks follow regarding the **exploration and exploitation tradeoff**. Consider  $k = 2$ :  $\Delta_1 = 0$ , and  $\Delta_2 = \Delta > 0$ . Then the above regret bound reads

$$\mathcal{R}_T^{\text{sto}}(\pi^{\text{ETC}}) \leq m\Delta + T\Delta \exp\left(-\frac{m\Delta^2}{4}\right) \quad (2.5)$$

Let us denote

$$m_0 = \lceil \frac{4}{\Delta^2} \log\left(\frac{T\Delta^2}{4}\right) \rceil \quad (2.6)$$

- **Gap dependent regret.**

- If  $m_0 \geq 1$ , then set  $m = m_0$

$$\mathcal{R}_T^{\text{sto}}(\pi^{\text{ETC}}) \leq \Delta + \frac{4}{\Delta} \left( \log\left(\frac{T\Delta^2}{4}\right) + 1 \right)$$

TL:  $\log(T)$  regret

- If  $m_0 < 1$ , then  $\frac{T\Delta^2}{4} < 1$  which implies that the gap is small  $\Delta < \frac{2}{\sqrt{T}}$ , in such case, set  $m = T/2$  will result in

$$\mathcal{R}_T^{\text{sto}}(\pi^{\text{ETC}}) = \frac{T}{2}\Delta \leq \sqrt{T}$$

- **Gap independent regret.** If  $m_0 \geq 1$ , then set  $m = m_0$  and note that

$$\frac{4}{\Delta} \left( \log\left(\frac{T\Delta^2}{4}\right) + 1 \right) \leq 2\sqrt{T} \sup_{x \geq 0} \frac{2 \log x + 1}{x} \leq 4\sqrt{e}\sqrt{T}$$

and thus

$$\mathcal{R}_T^{\text{sto}}(\pi^{\text{ETC}}) \leq \Delta + 4\sqrt{e}\sqrt{T}$$

This shows the worst case  $\sqrt{T}$  regret.

TL:  $\sqrt{T}$  regret

- **Gap independent policy.** So far the  $m$  depends on the knowledge of  $\Delta$ . Without knowing it and simply set  $m = T^{2/3}$ , we have

$$\begin{aligned} \mathcal{R}_T^{\text{sto}}(\pi^{\text{ETC}}) &\leq T^{2/3}\Delta + T^{2/3} \cdot T^{1/3}\Delta \exp\left(-\frac{T^{2/3}\Delta^2}{4}\right) \\ &\leq T^{2/3}\Delta + T^{2/3} \cdot 2 \sup_{x \geq 0} x \exp(-x^2) \asymp T^{2/3} \end{aligned}$$

TL:  $T^{2/3}$  regret

*Proof.* Observe the simple identity

$$\begin{aligned} \mathcal{R}_T^{\text{sto}}(\pi) &:= T \cdot \max_{i \in \mathcal{A}_k} \mu(i) - \mathbb{E} \left[ \sum_{t=1}^T \mu(u_t) \right] \\ &= \sum_{i=1}^k \Delta_i \mathbb{E}[N^T(i)] \end{aligned}$$

WLOG, assume  $\Delta_1 = 0$ , and  $\Delta_i > 0$  for  $i \geq 2$ . Let's control

$$\begin{aligned}
 \mathbb{E}[N^T(i)] &= \sum_{t=1}^T \mathbb{P}[u_t = i] \\
 &= m + \sum_{t=mk+1}^T \mathbb{P}[u_t = i] \\
 &= m + \sum_{t=mk+1}^T \mathbb{P}\left[\widehat{\mu}^{mk}(i) > \widehat{\mu}^{mk}(j), \forall j\right] \\
 &\leq m + \sum_{t=mk+1}^T \mathbb{P}\left[\widehat{\mu}^{mk}(i) > \widehat{\mu}^{mk}(1)\right] \\
 &\leq m + (T - mk) \mathbb{P}\left[\frac{1}{m} \sum_{z=0}^{m-1} (\mathbf{r}_{zk+i}(i) - \mathbf{r}_{zk+1}(1) - \mu(i) + \mu(1)) > \Delta_i\right] \\
 &\leq m + (T - mk) \exp\left(-\frac{m\Delta_i^2}{4}\right)
 \end{aligned}$$

Clearly  $B_z := \mathbf{r}_{zk+i}(i) - \mathbf{r}_{zk+1}(1) \in [-\sqrt{2}, \sqrt{2}]$  and that

$$\mathbb{P}\left[\frac{1}{m} \sum_{z=0}^{m-1} B_z - \mathbb{E}[B_z] > t\right] \leq \exp\left(-\frac{mt^2}{4}\right) \quad (2.7)$$

and the last line follows from Azuma-Hoeffding's inequality.  $\square$

## 2.2 Upper Confidence Bound (UCB) Algorithm

The next algorithm, UCB, is derived based on the *optimism principle*. The name **upper confidence bound** came from the Azuma Hoeffding's inequality

$$\mathbb{P}\left[\mu > \widehat{\mu} + \sqrt{\frac{2 \log(1/\delta)}{n}}\right] \leq \delta. \quad (2.8)$$

It shows an optimistic estimate of the true  $\mu$  based on  $n$ -empirical samples, with accuracy parameter  $\delta$ .

**Definition 4.** *Upper Confidence Bound (UCB) Algorithm* Define the

$$\text{UCB}^{t,\delta}(i) = \begin{cases} \infty & \text{if } N^t(i) = 0 \\ \widehat{\mu}^t(i) + \sqrt{\frac{2 \log(1/\delta)}{N^t(i)}} & \text{otherwise} \end{cases} \quad (2.9)$$

The UCB algorithm  $\pi^{\text{UCB}}$  implements the following

1. Input a small accuracy parameter  $\delta \in (0, 1)$
2. In round  $t$ , choose action

$$u_t := \arg \max_i \text{UCB}^{t,\delta}(i)$$

3. Observe the new reward and update the upper confidence bounds

**Theorem 2** (Regret for UCB Algorithm). Consider stochastic bandits with bounded rewards  $\mathbf{r}(i) \in [-1, 1]$  and  $\delta = \frac{1}{T(T+1)}$

$$\mathcal{R}_T^{\text{sto}}(\pi^{\text{UCB}}) \leq 2 \sum_{i=1}^k \Delta_i + \sum_{i:\Delta_i>0} \frac{16 \log(T+1)}{\Delta_i} \quad (2.10)$$

*Proof.* First, let us introduce one new notation.

$$\mathbf{R} = [r_{it}]_{i \in [k], t \in [T]} \in \mathbb{R}^{k \times T} \quad (2.11)$$

be the random reward matrix, where  $r_{i,t}, t \in [T]$  are i.i.d. draws from the distribution  $\mathcal{L}(\mathbf{r}(i))$ . We introduce the empirical mean as

$$\hat{\mu}_m(i) := \frac{1}{m} \sum_{t=1}^m r_{i,t}. \quad (2.12)$$

The Stochastic bandit model is stochastically equivalent to the following model: at each time  $t$ , if an arm  $i$  is pulled, then reveal the reward  $\mathbf{r}_t(i) = r_{i,N^t(i)}$ .

Again recall the basic identity,

$$\mathcal{R}_T^{\text{sto}}(\pi) := \sum_{i=1}^k \Delta_i \mathbb{E}[N^T(i)]$$

WLOG, assume  $\Delta_1 = 0$ , and  $\Delta_i > 0$  for  $i \geq 2$ . Let's control

$$\mathbb{E}[N^T(i)]$$

Choose an integer  $m_i$

$$m_i = \lceil \frac{8 \log(1/\delta)}{\Delta_i^2} \rceil < T$$

so to satisfy

$$\sqrt{\frac{2 \log(1/\delta)}{m_i}} < \frac{\Delta_i}{2} \quad (2.13)$$

Define two events

$$E_i = \left\{ \hat{\mu}_{m_i}(i) + \sqrt{\frac{2 \log(1/\delta)}{m_i}} < \mu(1) \right\} \quad (2.14)$$

$$F = \left\{ \forall t \in [T], \text{UCB}^{t,\delta}(1) > \mu(1) \right\} \quad (2.15)$$

**TL:** the distribution of  $\mathbf{r}_t(i)$  is the same as  $r_{i,1}$  though  $N^t(i)$  is a random variable.

**TL:**  $m_i$  conceptually is the right level of number of arms one need to pull to figure out in order to eliminate the bad arm.

Let's bound the probability of each event

$$\begin{aligned}
 \mathbb{P}[E_i^c] &= \mathbb{P}\left[\widehat{\mu}_{m_i}(i) + \sqrt{\frac{2\log(1/\delta)}{m_i}} \geq \mu(1)\right] \\
 &= \mathbb{P}\left[\widehat{\mu}_{m_i}(i) - \mu(i) \geq \Delta_i - \sqrt{\frac{2\log(1/\delta)}{m_i}}\right] \\
 &\leq \mathbb{P}\left[\widehat{\mu}_{m_i}(i) - \mu(i) \geq \frac{\Delta_i}{2}\right] \\
 &\leq \exp\left(-\frac{m_i\Delta_i^2}{8}\right) \leq \delta
 \end{aligned}$$

$$\begin{aligned}
 \mathbb{P}[F^c] &= \mathbb{P}\left[\exists t \in [T], \text{UCB}^{t,\delta}(1) \leq \mu(1)\right] \\
 &\leq \mathbb{P}\left[\exists m \in [T], \widehat{\mu}_m(1) + \sqrt{\frac{2\log(1/\delta)}{m}} \leq \mu(1)\right] \\
 &\leq T\delta
 \end{aligned}$$

Thus

$$\mathbb{P}[(E_i \cap F)^c] \leq \mathbb{P}[E_i^c] + \mathbb{P}[F^c] = (T+1)\delta. \quad (2.16)$$

On the event  $E_i \cap F$ , we claim that

$$N^T(i) \leq m_i. \quad (2.17)$$

If not, define

$$\tau := \inf\{t \in [T] : N^t(i) = m_i\} \leq T-1 \quad (2.18)$$

and at time  $\tau$ , arm  $i$  is pulled again. Then one must have

$$\text{UCB}^{\tau,\delta}(i) > \text{UCB}^{\tau,\delta}(1) \quad (2.19)$$

Recall that on the event  $E_i$ , we have

$$\text{UCB}^{\tau,\delta}(i) = \widehat{\mu}^\tau(i) + \sqrt{\frac{2\log(1/\delta)}{N^\tau(i)}} = \widehat{\mu}_{m_i}(i) + \sqrt{\frac{2\log(1/\delta)}{m_i}} < \mu(1) \quad (2.20)$$

yet on the event  $F$  we know

$$\text{UCB}^{\tau,\delta}(1) > \mu(1). \quad (2.21)$$

Therefore on the event  $E_i \cap F$ , (2.20) and (2.21) contradicts with (2.19).

Therefore, we have shown on the event  $E_i \cap F$ ,  $N^T(i) \leq m_i$ .

Now we are ready to bound

$$\begin{aligned}
 \mathbb{E}[N^T(i)] &= \mathbb{E}[N^T(i) \cdot \mathbf{1}_{E_i \cap F}] + \mathbb{E}[N^T(i) \cdot \mathbf{1}_{(E_i \cap F)^c}] \\
 &\leq m_i \cdot \mathbb{P}[E_i \cap F] + T \cdot \mathbb{P}[(E_i \cap F)^c] \\
 &\leq \frac{8\log(1/\delta)}{\Delta_i^2} + 1 + T(T+1)\delta
 \end{aligned}$$

Plug in  $\delta = \frac{1}{T(T+1)}$ , we know

$$\mathbb{E}[N^T(i)] \leq 2 + \frac{16 \log(T+1)}{\Delta_i^2} \quad (2.22)$$

and thus reach the final bound.  $\square$

### 3 Adversarial Bandits

Now we relieve the i.i.d. assumptions, and see that won't change the regret in the worst case  $\sqrt{T}$ . However, note in the best case, the regret in the stochastic bandit setting grows at a rate of  $\log(T)$ .

#### 3.1 Exponential Weight for Exploration and Exploitation (EXP3) Algorithm

**Definition 5 (EXP3 Algorithm).** Define the inverse propensity score estimate of the reward vector

$$\hat{\mathbf{r}}_t(i) = \begin{cases} \frac{\mathbf{r}_t(i)}{\mathbb{P}[u_t=i]} & \text{if } i = u_t \\ 0 & \text{otherwise} \end{cases} \quad (3.1)$$

At each round  $t$ , let  $u_t \sim \mathbf{p}_t$  where  $\mathbf{p}_t \in \mathbb{R}^k$  is a probability vector that sums up to 1.

1. Input a learning rate  $\eta$
2. In round  $t$ , choose action  $u_t \sim \mathbf{p}_t$  where  $\mathbf{p}_t \in \mathbb{R}^k$  is a probability distribution over arms
3. Observe the reward  $\mathbf{r}_t(u_t)$  scalar, and update

$$\mathbf{p}_{t+1}(i) = \frac{\mathbf{p}_t(i) \exp(\eta \hat{\mathbf{r}}_t(i))}{\sum_j \mathbf{p}_t(j) \exp(\eta \hat{\mathbf{r}}_t(j))} \quad (3.2)$$

**Theorem 3 (Regret for EXP3 Algorithm).** Consider the adversarial bandits with bounded rewards  $\mathbf{r}(i) \in [-1, 1]$  and  $\eta \in (0, 1)$

$$\mathcal{R}_T^{\text{adv}}(\pi^{\text{EXP3}}) \leq \frac{\log(k)}{\eta} + \eta \cdot Tk \quad (3.3)$$

**Remark.** Note this bound is optimized when  $\eta = \sqrt{\frac{\log(k)}{Tk}}$ , in this case

$$\mathcal{R}_T^{\text{adv}}(\pi^{\text{EXP3}}) \leq 2\sqrt{T \cdot k \log k}$$

*Proof.* We first introduce an inequality based on KL divergence  $D(\cdot \parallel \cdot)$ , define

$$\mathbf{p}_1(i) = \frac{\mathbf{p}_0(i) \exp(\eta \mathbf{r}(i))}{\sum_j \mathbf{p}_0(j) \exp(\eta \mathbf{r}(j))} \quad (3.4)$$

We then claim that  $\forall \mathbf{q}$

$$\langle \mathbf{r}, \mathbf{q} \rangle - \langle \mathbf{r}, \mathbf{p}_0 \rangle = \frac{D(\mathbf{q} \parallel \mathbf{p}_0) - D(\mathbf{q} \parallel \mathbf{p}_1)}{\eta} + \frac{D(\mathbf{p}_0 \parallel \mathbf{p}_1)}{\eta}. \quad (3.5)$$

To derive this, notice

$$D(\mathbf{q} \parallel \mathbf{p}) = \langle \mathbf{q}, \log \mathbf{q} - \log \mathbf{p} \rangle \quad (3.6)$$

and thus the RHS equals

$$\begin{aligned} & \frac{\langle \mathbf{q}, \log \mathbf{q} - \log \mathbf{p}_0 \rangle - \langle \mathbf{q}, \log \mathbf{q} - \log \mathbf{p}_1 \rangle}{\eta} + \frac{\langle \mathbf{p}_0, \log \mathbf{p}_0 - \log \mathbf{p}_1 \rangle}{\eta} \\ &= \frac{1}{\eta} (\langle \mathbf{q}, \log \mathbf{p}_1 - \log \mathbf{p}_0 \rangle - \langle \mathbf{p}_0, \log \mathbf{p}_1 - \log \mathbf{p}_0 \rangle) \\ &= \frac{1}{\eta} (\langle \mathbf{q}, \mathbf{r} \rangle - \langle \mathbf{p}_0, \mathbf{r} \rangle + \text{const.} \cdot \langle \mathbf{q} - \mathbf{p}_0, \mathbf{1} \rangle) \\ & \text{(notice } \log \mathbf{p}_1 - \log \mathbf{p}_0 = \eta \mathbf{r} + \text{const.} \cdot \mathbf{1}) \end{aligned}$$

The claim is thus proved.

A second fact we will use is a bound for the KL divergence through the local norm: suppose  $z \in [-1, 1]$ , then we have

$$\exp(z) - 1 - z \leq z^2 \quad (3.7)$$

and thus

$$\frac{D(\mathbf{p}_0 \parallel \mathbf{p}_1)}{\eta} \leq \eta \sum_{i=1}^k \mathbf{p}_0(i) \mathbf{r}^2(i). \quad (3.8)$$

The proof is due to the fact

$$\begin{aligned} D(\mathbf{p}_0 \parallel \mathbf{p}_1) &= \log(\langle \mathbf{p}_0, \exp(\eta \mathbf{r}) \rangle) - \langle \mathbf{p}_0, \eta \mathbf{r} \rangle \\ &\leq \langle \mathbf{p}_0, \exp(\eta \mathbf{r}) - 1 \rangle - \langle \mathbf{p}_0, \eta \mathbf{r} \rangle \quad \text{notice } \log(1+z) \leq z, \forall z \geq -1. \\ &= \sum_{i=1}^k \mathbf{p}_0(i) (\exp(\eta \mathbf{r}(i)) - 1 - \eta \mathbf{r}(i)) \quad \text{notice } \exp(z) - 1 - z \leq z^2, \forall z \in [-1, 1] \\ &\leq \eta^2 \sum_{i=1}^k \mathbf{p}_0(i) \mathbf{r}^2(i). \end{aligned}$$

So far we have derived

$$\langle \mathbf{r}, \mathbf{q} \rangle - \langle \mathbf{r}, \mathbf{p}_0 \rangle \leq \frac{D(\mathbf{q} \parallel \mathbf{p}_0) - D(\mathbf{q} \parallel \mathbf{p}_1)}{\eta} + \eta \sum_{i=1}^k \mathbf{p}_0(i) \mathbf{r}^2(i). \quad (3.9)$$

Recursively using the above (3.9), we can obtain the following telescoping inequality for the EXP3 algorithm

$$\langle \hat{\mathbf{r}}_t, \mathbf{q} \rangle - \langle \hat{\mathbf{r}}_t, \mathbf{p}_t \rangle \leq \frac{D(\mathbf{q} \parallel \mathbf{p}_t) - D(\mathbf{q} \parallel \mathbf{p}_{t+1})}{\eta} + \eta \sum_{i=1}^k \mathbf{p}_t(i) \hat{\mathbf{r}}_t^2(i). \quad (3.10)$$



Notice that due to the inverse propensity rule being unbiased, we have

$$\mathbb{E}_{u_t}[\langle \widehat{\mathbf{r}}_t, \mathbf{q} \rangle] = \langle \mathbb{E}_{u_t}[\widehat{\mathbf{r}}_t], \mathbf{q} \rangle = \langle \mathbf{r}_t, \mathbf{q} \rangle \quad (3.11)$$

$$\mathbb{E}_{u_t}[\langle \widehat{\mathbf{r}}_t, \mathbf{p}_t \rangle] = \langle \mathbb{E}_{u_t}[\widehat{\mathbf{r}}_t], \mathbf{p}_t \rangle = \langle \mathbf{r}_t, \mathbf{p}_t \rangle = \mathbb{E}_{u_t}[\mathbf{r}_t(u_t)] \quad (3.12)$$

and thus

$$\langle \mathbf{r}_t, \mathbf{q} \rangle - \mathbb{E}_{u_t}[\mathbf{r}_t(u_t)] \leq \mathbb{E}_{u_t} \left[ \frac{D(\mathbf{q} \parallel \mathbf{p}_t) - D(\mathbf{q} \parallel \mathbf{p}_{t+1})}{\eta} \right] + \eta \sum_{i=1}^k \mathbf{p}_t(i) \mathbb{E}_{u_t}[\widehat{\mathbf{r}}^2(i)] \quad (3.13)$$

where

$$\mathbb{E}_{u_t}[\widehat{\mathbf{r}}^2(i)] = \frac{\mathbf{r}_t^2(i)}{\mathbf{p}_t(i)} \quad (3.14)$$

and thus

$$\langle \mathbf{r}_t, \mathbf{q} \rangle - \mathbb{E}_{u_t}[\mathbf{r}_t(u_t)] \leq \mathbb{E}_{u_t} \left[ \frac{D(\mathbf{q} \parallel \mathbf{p}_t) - D(\mathbf{q} \parallel \mathbf{p}_{t+1})}{\eta} \right] + \eta \sum_{i=1}^k \mathbf{p}_t(i) \frac{\mathbf{r}_t^2(i)}{\mathbf{p}_t(i)} \quad (3.15)$$

$$\leq \mathbb{E}_{u_t} \left[ \frac{D(\mathbf{q} \parallel \mathbf{p}_t) - D(\mathbf{q} \parallel \mathbf{p}_{t+1})}{\eta} \right] + \eta k \quad (3.16)$$

and thus marginally we have

$$\langle \mathbf{r}_t, \mathbf{q} \rangle - \mathbb{E}[\mathbf{r}_t(u_t)] \leq \mathbb{E} \left[ \frac{D(\mathbf{q} \parallel \mathbf{p}_t) - D(\mathbf{q} \parallel \mathbf{p}_{t+1})}{\eta} \right] + \eta \sum_{i=1}^k \mathbf{p}_t(i) \frac{\mathbf{r}_t^2(i)}{\mathbf{p}_t(i)} \quad (3.17)$$

$$\leq \mathbb{E} \left[ \frac{D(\mathbf{q} \parallel \mathbf{p}_t) - D(\mathbf{q} \parallel \mathbf{p}_{t+1})}{\eta} \right] + \eta k \quad (3.18)$$

summing over  $t = 1, 2, \dots, T$

$$\begin{aligned} \mathcal{R}_T^{\text{adv}}(\pi^{\text{EXP3}}) &\leq \frac{D(\mathbf{q} \parallel \mathbf{p}_1) - \mathbb{E}[D(\mathbf{q} \parallel \mathbf{p}_{T+1})]}{\eta} + T \cdot \eta k \\ &\leq \frac{\log(k)}{\eta} + \eta \cdot Tk \end{aligned}$$

□

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